Course: Automata Theory

Lecture 5: Kleene's Theorem, Symbolic Logic and Formal Systems

Lecturer: Martha Gichuki

Course description

- The course begins with an introduction to logic and formal grammar where learners will do a recap on sets, logic and truth tables, sequences, relations and functions
- A coverage of finite state machines, Push Down automata and Turing Machines (The Church's thesis) will culminate the study of various models of computation.
- Formal language and grammar will then follow to enable learners differentiate regular and context free languages.
- An evaluation of the computability and complexity of practical computational problems which are the foundations of automata theory will then be done and the outcome will be problem description.

Learning outcomes:

Lecture 5: Kleene's Theorem, Symbolic Logic and Formal Systems

At the end of the lecture you will be able to:

i. Define Kleene's Theorem.
ii. Describe Formal Systems.
iii. Differentiate symbolic logic and formal systems

5.1 Kleene's Theorem

Introduction:

- From lecture one, we indicated that Automata Theory is closely related to **formal language theory** because automatons are often classified by the class of formal languages they can recognize.

Introduction…

- Recall from lecture 1 that an automaton is anchored on the basic concepts of **symbols** (characters)**, words** (symbols) **alphabets** (set of symbols and letters) and **strings.**
- A **Language** is a set of words formed by symbols in each alphabet and according to Kleene ***Closure,*** a language may be thought of as a sub-set of possible words.

Introduction….

- The set of all possible words may in turn be thought of as the set of all possible concatenations of strings denoted as * and the super script $^{(*)}$ is called the **Kleene Star**.
- An automaton has a mechanism to **read inputs** which is a **string over a given alphabet**

Kleene’s Theorem

- Kleene’s theorem states that Finite Automatons accept regular languages and regular expressions are used to describe regular languages.
- A set of regular is defined over an alphabet Σ recursively and all the elements of that set is from regular expressions

5.2 Symbolic Logic

- Symbolic logic provides a way to express logical expressions by using symbols and variables instead of using natural languages such as English.
- This kind of standard expression helps remove vagueness which would otherwise occur if we used natural languages

What are logical expressions?

- Recall from Lecture 3 that logical statements are statements with a truth value: either true or false however, there is no need of knowing whether a logical expression is true or false, all we need to know is that it has a truth value.
- Consider a question like 'What are you eating?' This has no truth value.
- Consider a statement like “It is raining” This expression is either true or false.

Propositions:

- The smallest logical expression possible, which, if broken down further loses meaning is called a **proposition**.

Example 1 of a proposition:

- *'James and John live together'* cannot be broken down without a loss in meaning. *'James lives together'* does not make sense.

Example 2 of a proposition:

- '*Joyce and Joy go to work*" can be broken down into *'Joyce goes to work'* and *'Joy goes to work,'*
- There is no evidence showing that Joyce and Joy go to work at the same time.

Propositions in Symbolic Logic

- In symbolic logic, propositions may be represented by capital letters such as A or B, or lower-case letters such as p, q, r.
- Example '*My house is big*' may represented by A, and '*My puppy is white*' may be represented by B.
- A = '*My house is big*'
- B = '*My puppy is white*'

Writing Propositions

- Propositions are written in the affirmative.
- Example we use the not symbol (~) to make a **negation** (a not statement) instead of writing '*my puppy is not white*',
- We write using symbols, as follows: ~ B. Therefore, logically, negation changes an expression's truth value

Writing Propositions

***Example**: Consider the proposition* 'My puppy is white', then B would be True, and ~ B would be False,

- B = '*My puppy is white*' (*True*)
- ~ B = '*My puppy is not white*' (*False*)

Writing Propositions

- ***Example:*** *Consider the proposition* 'My house is small', then A would be false, and ~ A would be true.
- A = '*My house is small*' (*False*)
- ~ A= '*My house is big*' (*True*)

Truth Tables

- From Lecture 3 recall that a **truth table** lists whether something is true or false and as stated above, the act of negation changes all the true values into false and the false values to true as shown below.

P	~ P
True	False
False	True

- As seen in Lecture 3 & 4, to construct a truth table:

i. A column for every proposition in the expression must be created,

ii. A row for every possible true and false value is created

This ensures that all the possible combinations of a given set of propositions are catered for.

- Example: For two propositions, a truth table for all the possible truth value combinations will have four rows for T/F possibilities and two columns for the propositions. In general, there will be **2^n rows for *n* different propositions**.

P	Q
True	True
True	False
False	True
False	False

Logical operators

- Logical operators are needed to link together propositions and build more complex logical expressions.

- They operate in the same way as +, -, ×, and ÷ operators used to link mathematical expressions.

- The basic logical operators, along with **negation (∼)**, are **conjunction (Λ)**, **disjunction (V)**, **conditional (→)**, and **biconditional (↔)** operators

5.3 Formal Systems

- Formal systems are also known as **logistic systems** in abstract terms

- Formal systems have a ***formal language comprising of symbols*** acted on by ***certain rules*** allowed in the system.

Formal Systems…

- The rules are developed from axioms therefore formal systems have inbuilt formulas that use finite combinations of primitive symbols.

- Model structure are used to interpret the symbols of a formal system and often used in conjunction with formal systems.

What constitutes formal systems?

A formal system has the following components:

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i. A **finite alphabet** of symbols; *finite* because we need precise models.
ii. A **syntax** that defines which strings of symbol are in the language of the formal system.
iii. A **decidable** set of **axioms** and a finite set of **rules** from which the set of **theorems** of the system is generated and the rules must take a finite number of steps to apply.

1. A finite alphabet of symbols, finite because we need precise models.

- **Example:** Consider a machine, $\mathbf{M_2}$ which accepts the language: {W∈{a,b}*: W has an odd number of ***a's*** and an even number of ***b's***}

- *The finite alphabet symbols are **a** and **b***

2. A **syntax** - which strings of the symbols are in the language of the formal system?

- **Example:** Consider a machine, $\mathbf{M_2}$ which accepts the language: {W∈{a,b}*: W has an odd number of ***a's*** and an even number of ***b's***}
- *The syntax here is how the accepted strings will be concatenated i.e.* ***odd number of a's*** *and an* ***even number of b's***
- *E.g. possible strings in W could be* ***{{abb}, {aaabb}, {aaaaabbbb}, {aaaaaaaabbbb}, {aaaaaaaaaaabbbbbb}, …}***

3. A **decidable** set of **axioms** and a finite set of **rules** from which the set of **theorems** of the system is generated.

- **Example:** Consider a machine, $\mathbf{M_2}$, which accepts the language: {W∈{a,b}*: W has an odd number of ***a's*** and an even number of ***b's***}
- *The rules are {W∈{a,b}*: W has an odd number of a's and an even number of b's}.*
- *Strings outside {**a,b**} cannot be accepted and the concatenation rule must be followed*

Review Questions

1. For any alphabet $\sum$, the set of all strings over $\sum$ is denoted by $\sum^*$, suppose $\sum$ is in the two-symbol alphabet {a, b}, what are the possible subsets of $\sum^*$?

- *Possible subsets= {Ø, {a}, {b}, {ab}, {ba}, {aa}, {bb}, {aba}, {bab}, {abba}, {baab}, {bbaa}, {aabb}, ...}*

2. ***Describe the following sets by regular expressions***:

i. {01, 10}

Set of alternate zeros and ones in pairs

ii. {101}

Set of triple palindrome zeros and ones

iii. {λ,1,1,1,1,1,1…}

Set of single digits of one with λ included

iv. {1,1,1,1,1,1,1…}

Set of single digits of one

3. If C is a set with *c* elements, how many elements are in the *power set of C*? Explain your answer.

Power set of C has $2^{|C|}$ *where* |C| is the cardinality (length) of C

References

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